

# Bayesian equilibria in Bayesian games

## Sealed bid auction

Two players, both willing to buy an object. Their values and bids lie in  $[0, 1]$

allocation function:  $O_1(b_1, b_2) = I\{b_1 \geq b_2\}$ ;  $O_2(b_1, b_2) = I\{b_2 > b_1\}$

beliefs:  $f(\theta_2 | \theta_1) = 1, \forall \theta_1, \theta_2$        $f(\theta_1, \theta_2) = 1, \forall (\theta_1, \theta_2) \in [0, 1]^2$   
 $f(\theta_1 | \theta_2) = 1, \forall \theta_1, \theta_2$

① First price auction: if  $b_1 \geq b_2$ , player 1 wins and pays her bid  
ow, player 2 wins and pays her bid

$$u_1(b_1, b_2, \theta_1, \theta_2) = (\theta_1 - b_1) I\{b_1 \geq b_2\}$$

$$u_2(b_1, b_2, \theta_1, \theta_2) = (\theta_2 - b_2) I\{b_1 < b_2\}$$

$b_1 = s_1(\theta_1), b_2 = s_2(\theta_2)$ , assume  $s_i(\theta_i) = \alpha_i \theta_i, \alpha_i > 0, i=1, 2$ .

To find the BE, we need to find the  $s_i^*$  (or  $\alpha_i^*$ ) that maximizes the ex-interim utility of player  $i$

$$\max_{\sigma_i} u_i(\sigma_i, \sigma_i^* | \theta_i),$$

For player 1, this reduces to:

$$\max_{b_1 \in [0, \alpha_2]} \int_0^1 f(\theta_2 | \theta_1) (\theta_1 - b_1) I\{b_1 \geq \alpha_2 \theta_2\} d\theta_2 \quad \left( \begin{array}{l} \text{since } \theta_2 \in [0, 1] \\ b_1 \text{ never needs} \\ \text{to be larger than} \\ \alpha_2 \end{array} \right)$$
$$= \max_{b_1 \in [0, \alpha_2]} (\theta_1 - b_1) \frac{b_1}{\alpha_2} \Rightarrow b_1^* = \begin{cases} \frac{\theta_1}{2} & \text{if } \alpha_2 > \frac{\theta_1}{2} \\ \alpha_2 & \text{ow} \end{cases}$$

$$s_1^*(\theta_1) = \min \left\{ \frac{\theta_1}{2}, \alpha_2 \right\}, s_2^*(\theta_2) = \min \left\{ \frac{\theta_2}{2}, \alpha_1 \right\}$$

if  $\alpha_1 = \alpha_2 = 1/2$ , then  $(\frac{\theta_1}{2}, \frac{\theta_2}{2})$  is a BE

In the Bayesian game induced by uniform prior on first price auction, bidding half the true value is a Bayesian equilibrium.

(2) Second price auction: highest bidder wins but pays the second highest bid.

$$u_1(b_1, b_2, \theta_1, \theta_2) = (\theta_1 - b_2) I\{b_1 \geq b_2\}$$

$$u_2(b_1, b_2, \theta_1, \theta_2) = (\theta_2 - b_1) I\{b_1 < b_2\}$$

Player 1's bidding problem is to maximize

$$\begin{aligned} & \int_0^1 f(\theta_2 | \theta_1) (\theta_1 - s_2(\theta_2)) I(b_1 \geq s_2(\theta_2)) d\theta_2 \\ &= \int_0^1 1 \cdot (\theta_1 - \alpha_2 \theta_2) I(\theta_2 \leq \frac{b_1}{\alpha_2}) d\theta_2 \\ &= \frac{1}{\alpha_2} (b_1 \theta_1 - \frac{\theta_1^2}{2}) \Rightarrow \text{maximized when } b_1 = \theta_1 \end{aligned}$$

Similarly for  $b_2 = \theta_2$ .

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If the distributions of  $\theta_1$  and  $\theta_2$  were arbitrary but independent the maximization problem would have been

$$\int_0^{b_1/\alpha_2} f(\theta_2) (\theta_1 - \alpha_2 \theta_2) d\theta_2 = \theta_1 F\left(\frac{b_1}{\alpha_2}\right) - \alpha_2 \int_0^{b_1/\alpha_2} \theta_2 f(\theta_2) d\theta_2$$

differentiating wrt  $b_1$ , we get

$$\theta_1 \frac{1}{\alpha_2} f\left(\frac{b_1}{\alpha_2}\right) - \alpha_2 \cdot \frac{b_1}{\alpha_2} f\left(\frac{b_1}{\alpha_2}\right) \cdot \frac{1}{\alpha_2} = 0$$

$$\Rightarrow \frac{1}{\alpha_2} f\left(\frac{b_1}{\alpha_2}\right) (b_1 - \theta_1) = 0 \Rightarrow b_1 = \theta_1, (\text{similar for } 2), \text{ if } f\left(\frac{b_1}{\alpha_2}\right) > 0$$

For any independent, positive prior, bidding true type is a BE of the induced Bayesian game in second price auction.