

Claim: Suppose f satisfies SP, ONTO, ANON, then

$$f(P) = \text{median}(p_1, \dots, p_n, y_1, \dots, y_{n-1}).$$

Case 1: a is a phantom peak - proved

Case 2: a is an agent peak

We will prove this for 2 players. The general case repeats this argument.

Claim: $N = \{1, 2\}$, let P and P' be such that

$$P_i(1) = P'_i(1), \forall i \in N. \text{ Then } f(P) = f(P').$$

Proof: Let $a = P_1(1) = P'_1(1)$, and $P_2(1) = P'_2(1) = b$.

$$f(P) = x \text{ and } f(P', P_2) = y$$

Since f is SP, $x P_i y$ and $y P'_i x$

Since peaks of P and P' are the same, if x, y are on the same side of the peak, they must be the same, as the domain is single peaked.

The only other possibility is that x and y fall on different sides of the peak. We show that this is impossible.

WLOG $x < a < y$ and $a < b$

f is SP+ONTO $\Leftrightarrow f$ is SP+PE

PE requires $f(P) \in [a, b]$, but $f(P) = x < a \rightarrow$

now repeat this argument for $(P'_1, P_2) \rightarrow (P'_1, P'_2)$ \square

Profile: $(P_1, P_2) = P$, $P_1(1) = a$, $P_2(1) = b$

y_1 is the phantom peak.

by assumption, median (a, b, y_1) is an agent peak

WLOG assume the median is a .

Assume for contradiction $f(P) = c \neq a$.

By PE, c must be within a and b . We have two cases to consider: $b < a < y_1$ and $y_1 < a < b$.

Case 2.1: $b < a < y_1$, by PE $c < a$

construct P'_1 s.t. $P'_1(1) = a = P_1(1)$

and y_1, P'_1, c (possible since they are on different sides of a)

by the earlier claim, $f(P) = c \Rightarrow f(P'_1, P_2) = c$.

now consider the profile (P'_1, P_2)

peak at the rightmost

$P_2(1) = b < y_1 \leq P'_1(1)$, hence the median of $\{b, y_1, P'_1(1)\}$ is y_1 (which is a phantom peak, hence case 1 applies).

$$f(P'_1, P_2) = y_1.$$

But $y_1, P'_1 \in c$ (by construction) and $f(P'_1, P_2) = c$
agent 1 manipulates $P'_1 \rightarrow P'_1$, contradiction to f being SP.

case 2.2: $y_1 < a < b$, PE $\Rightarrow a < c$

construct P'_1 s.t. $P'_1(1) = a = P_1(1)$ and $y_1, P'_1 \in c$
 $f(P'_1, P_2) = c$ (by claim)

consider (P'_1, P_2) , $P'_1(1) \leq y_1 < b \Rightarrow f(P'_1, P_2) = y_1$
but $y_1, P'_1 \in c$, hence manipulable by agent 1.

This completes the proof for two agents (case 2). For the generalization to n players, see Moulin (1980)

"On strategyproofness and single peakedness".